

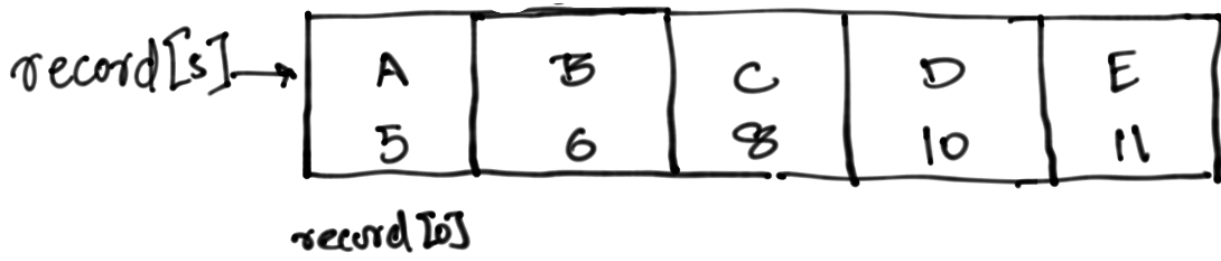
## Data - Structure Question

(1) Assume that you are the owner of a company that has 5 employees.

(2) Each employee is represented by  
{ name, salary }

Which data-structure (that you know) will you use to represent these 5 employees?

# Arrays



# Arrays

record[5] →

A	B	C	D	E
5	6	8	10	11

record[i]

Suppose a new employee [F, 12] joins the company. How will you update your data-structure?

# Arrays

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Suppose a new employee {F, 10} joins the company. How will you update your data-structure?

△ Add {F, 10} @ record[6]

# Arrays

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5	6	8	10	11

record[i]

Suppose a new employee {F, 10} joins the company. How will you update your data-structure?

△ Add {F, 10} @ record[6]

😞 But record[6] does not exist. The array record is of size 5 only.

Now what will you do?

(1) Make a new array `newrecord` of size 6

(2) Copy the five records from `record` to `newrecord`.

(3) Add `{F, 10}` @ `newrecord[6]`

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(2) Copy the five records from `record` to `newrecord`.

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Q: What is the running time of this procedure?

A:  $O(n)$  if there are  $n$  records in record array.



## When an employee leaves

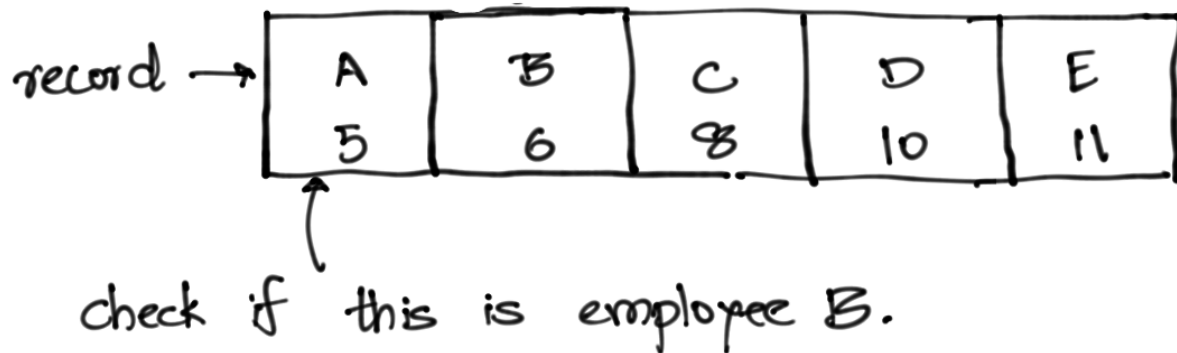
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record →

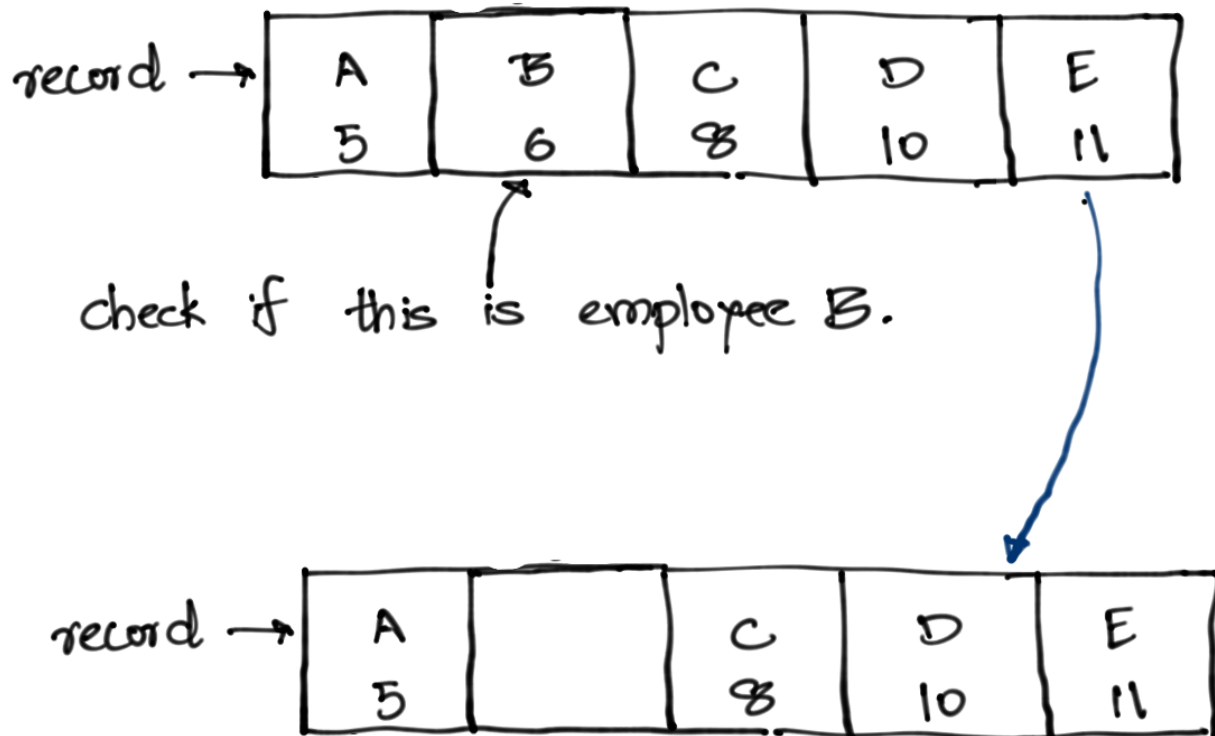
A	B	C	D	E
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check if this is employee B.

## When an employee leaves

Suppose employee B leaves, so we have to remove the record of employee B.

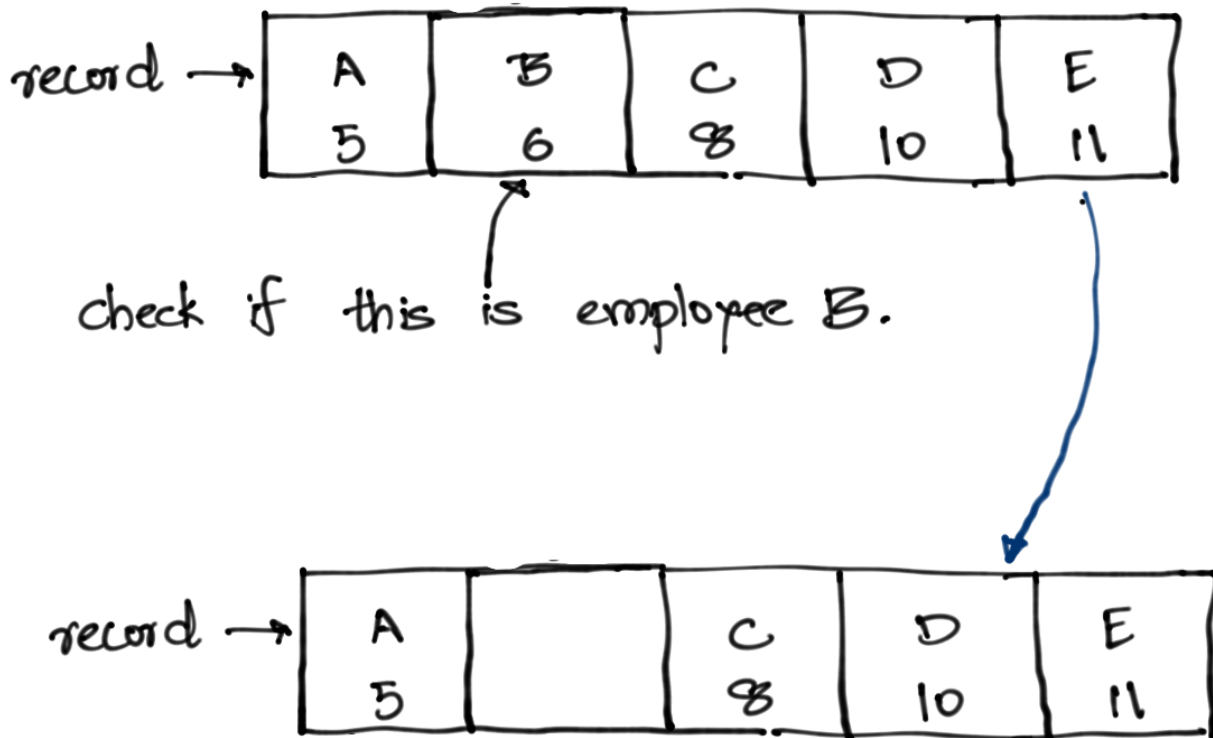
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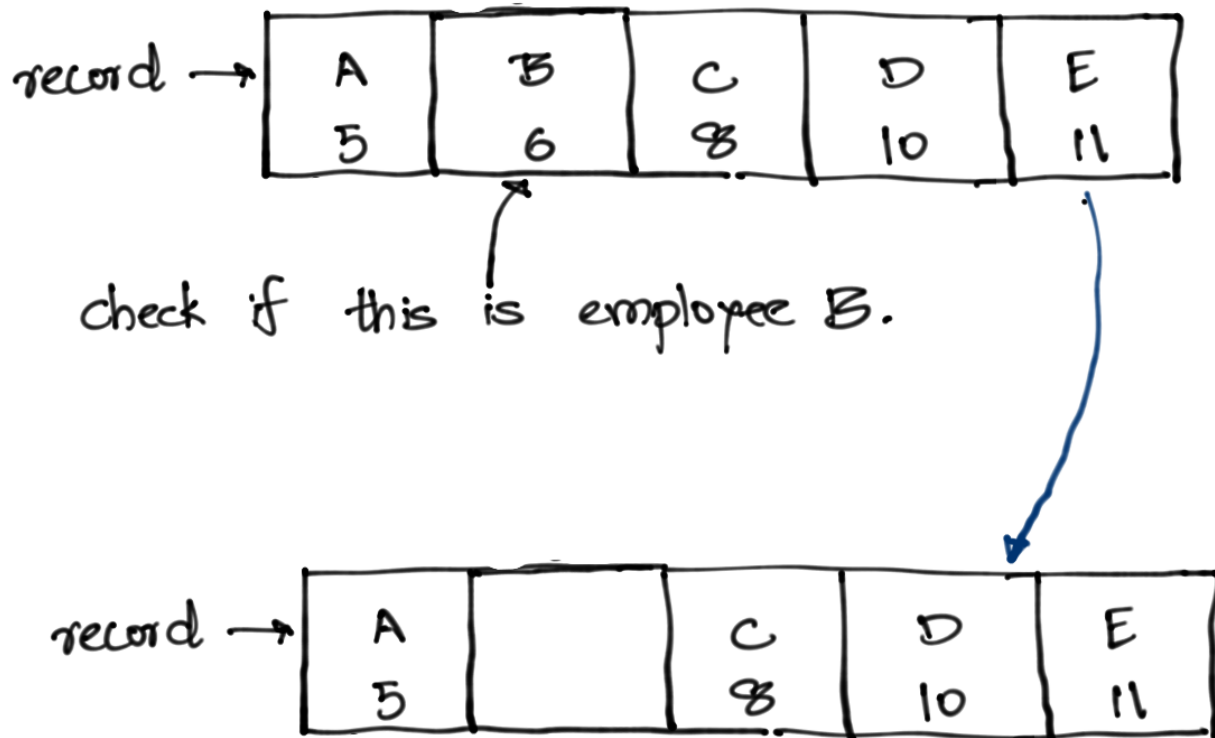


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## Holy Grail for data-structure

The space taken by your data-structure should be proportional to the number of current employees in the company.

# Performance of Arrays.

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Deletion	$O(n)$
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A: Give me the 5<sup>th</sup> element of record array.

return record[5]

Running Time = ??

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So, array nearly always give worst case performance.

Q: Where are arrays good!

A: Give me the 5<sup>th</sup> element of record array.

return record[5]

Running Time =  $O(1)$ .

## Performance of Arrays.

Insert	$O(n)$
Deletion	$O(n)$
Search	$O(n)$
Get the $k^{\text{th}}$ element	$O(1)$ .



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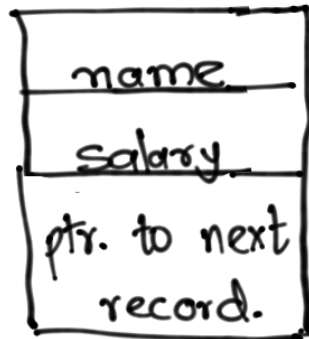
Mimics arrays.

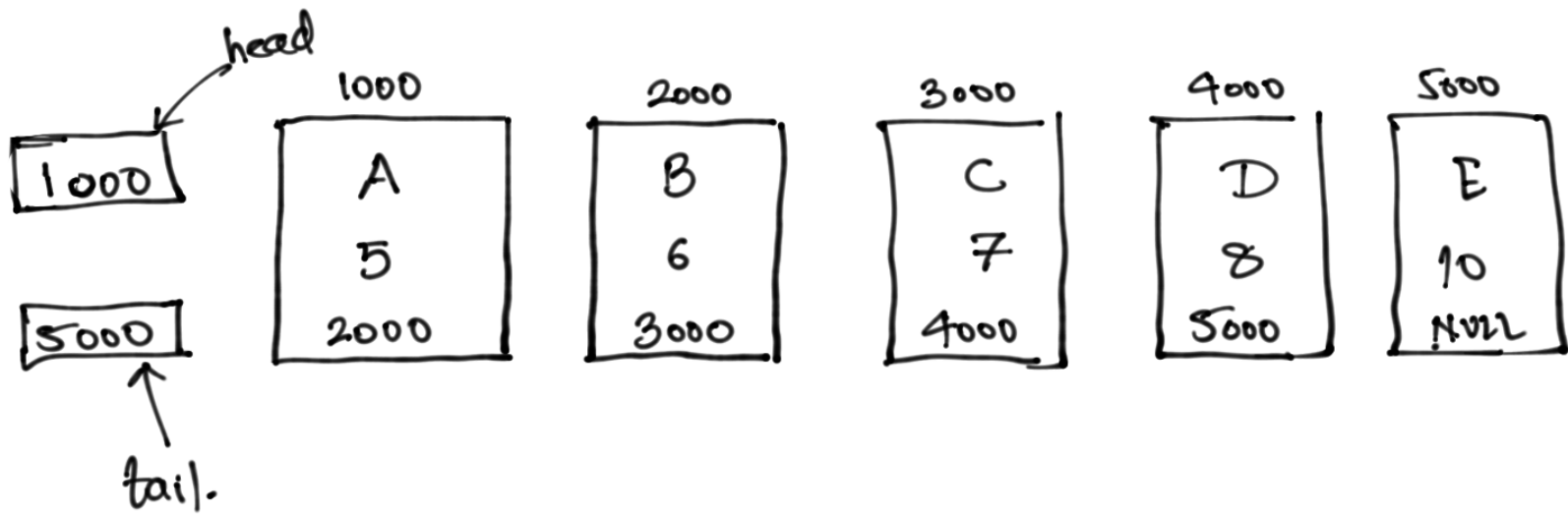
# Holy Grail for data-structure

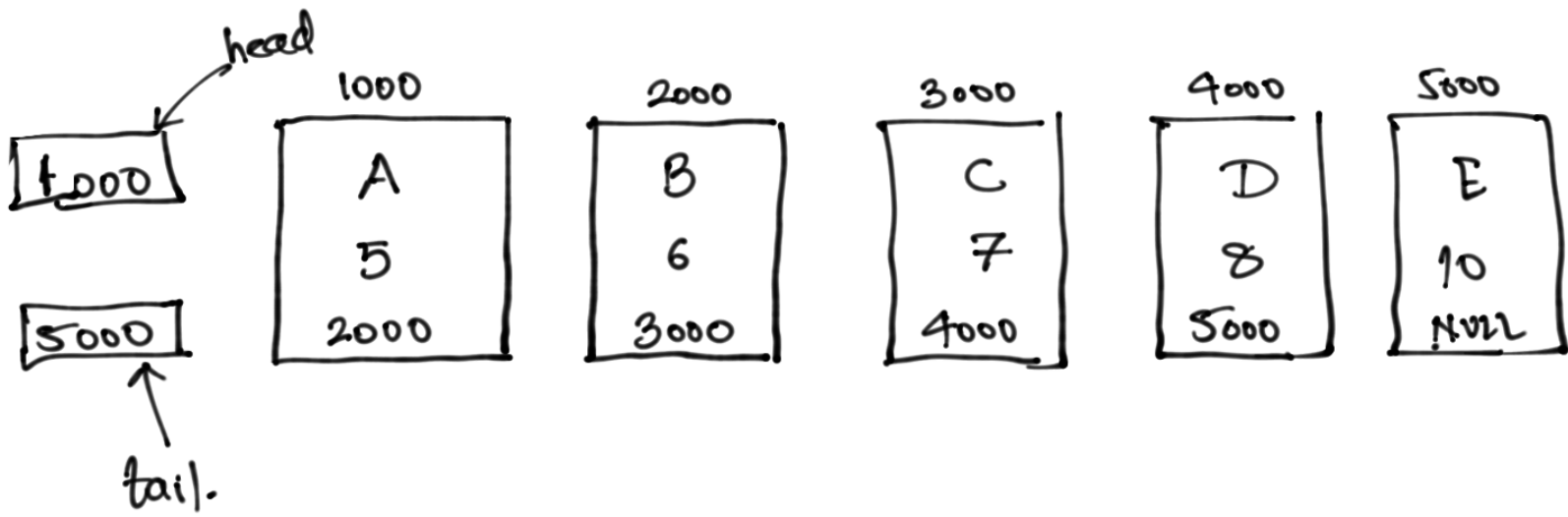
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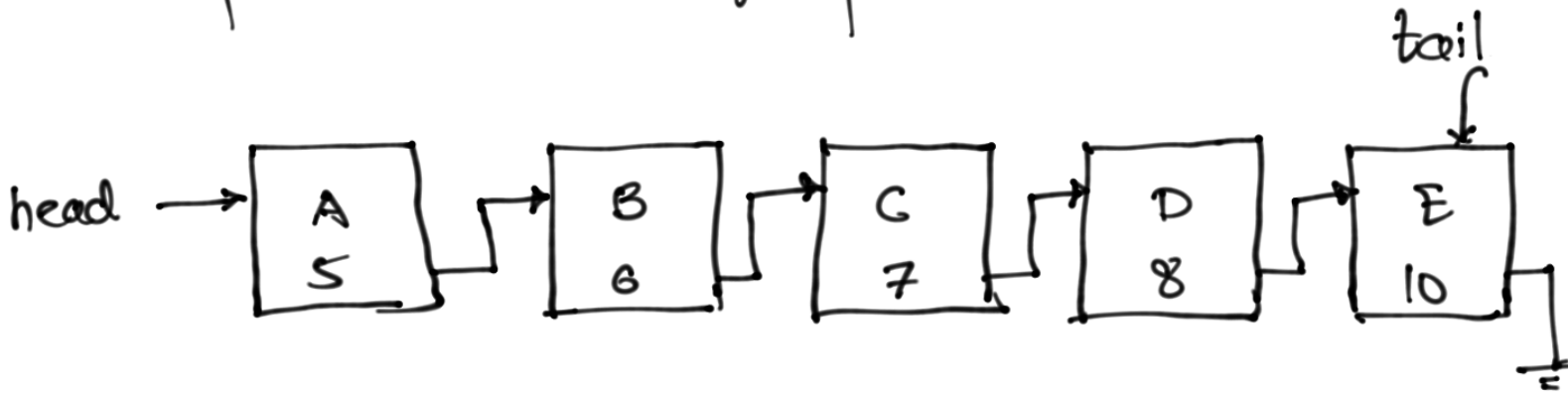
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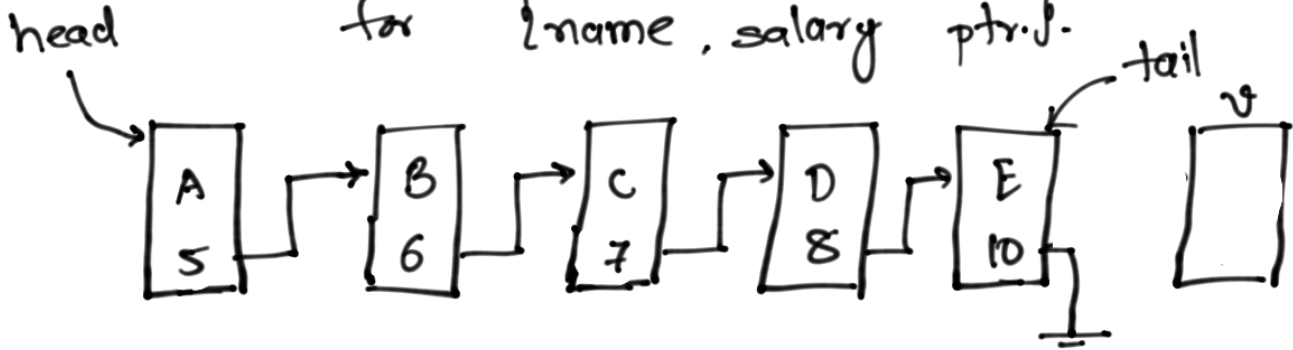
However, we will have a simpler pictorial view of pointers.



Insert  $\{F, 10\}$ .

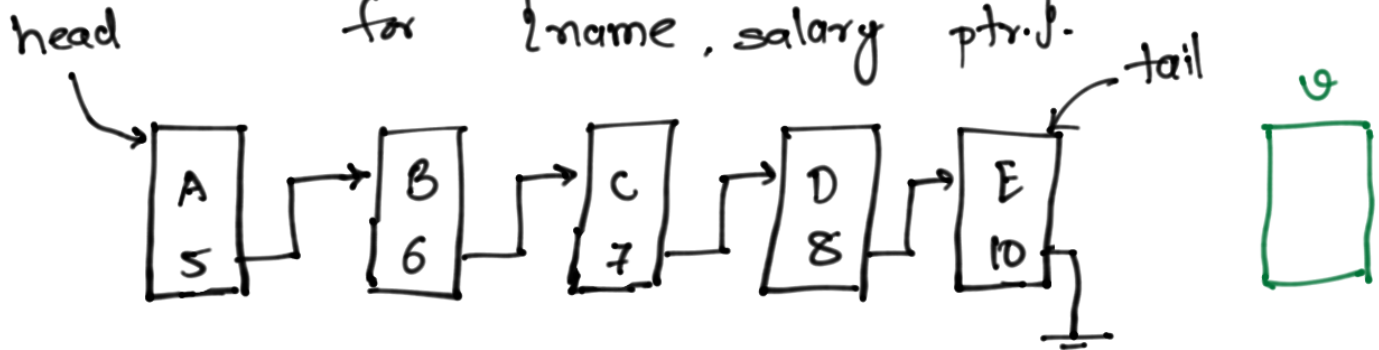
Insert {F, 10}.

①  $v \leftarrow$  Request the OS to allocate enough space for {name, salary ptr.}.

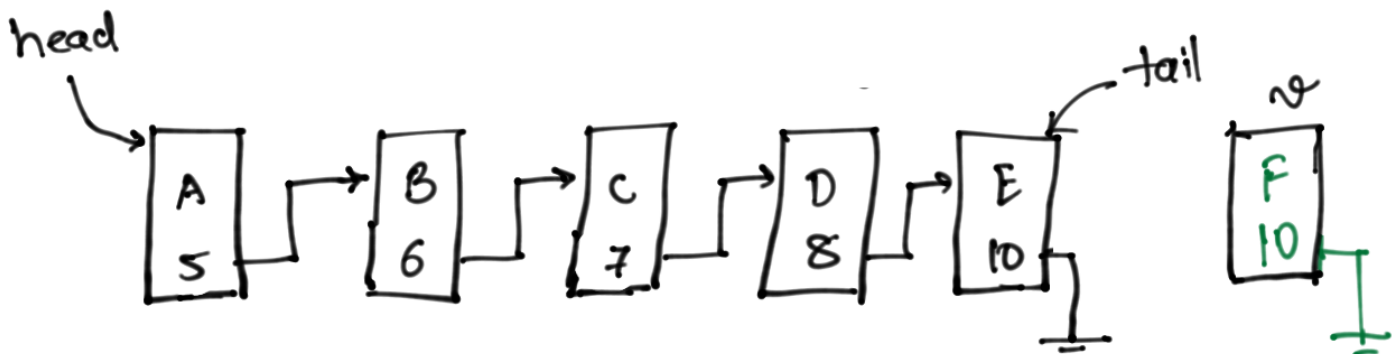


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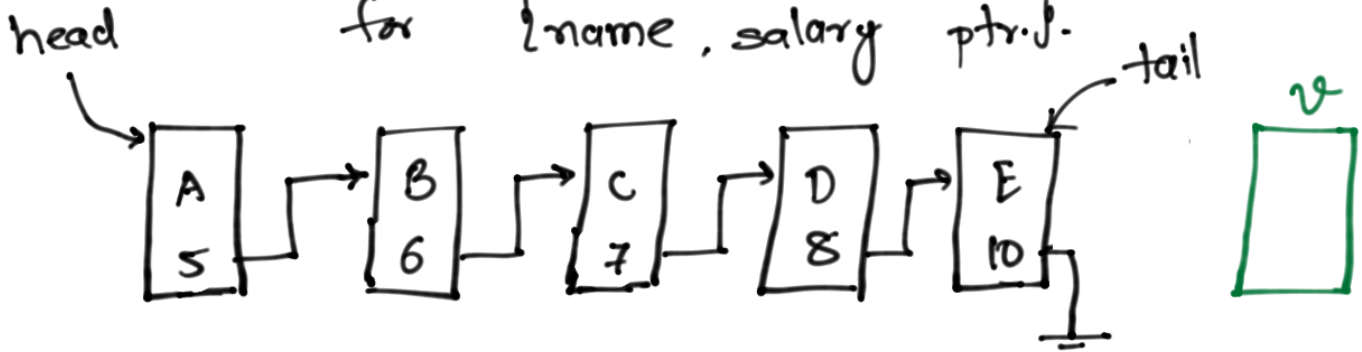
②  $v.name \leftarrow F$  ;  
 $v.salary \leftarrow 10$  ;  
 $v.ptr \leftarrow null$



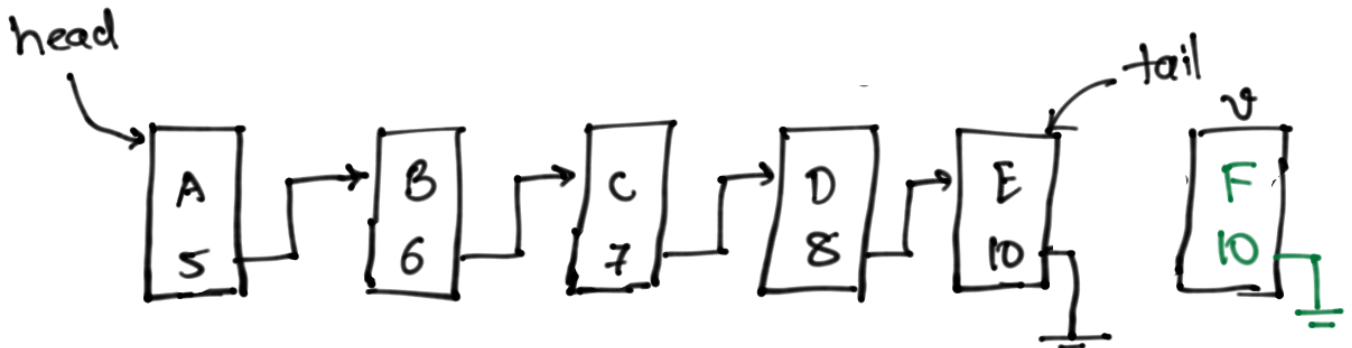


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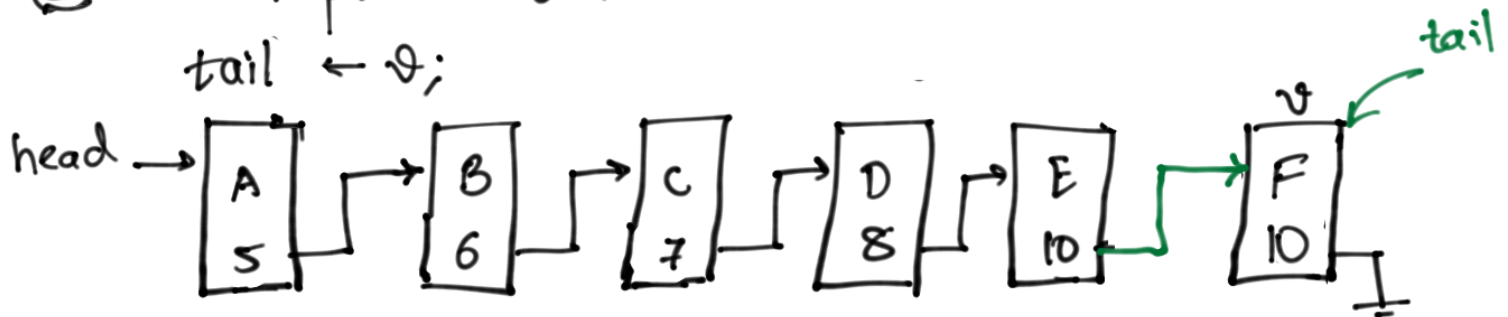
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②  $v.name \leftarrow F$  ;  
 $v.salary \leftarrow 10$  ;  
 $v.ptr \leftarrow NULL$



③  $tail.ptr \leftarrow v$  ;  
 $tail \leftarrow v$  ;



```
Insert (name, salary)
{
```

```
    v ← allocate me a new memory location;
```

```
    v.name ← name;
```

```
    v.salary ← salary;
```

```
    v.ptr ← NULL;
```

```
    if (head is NULL).
```

```
    {
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```
    if (head is NULL). // list is empty
```

```
    {
```

```
        head ← v;
```

```
        tail ← v;
```

```
    }
```

```
    else
```

```
    {
```

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Q: What is the running time of this procedure!

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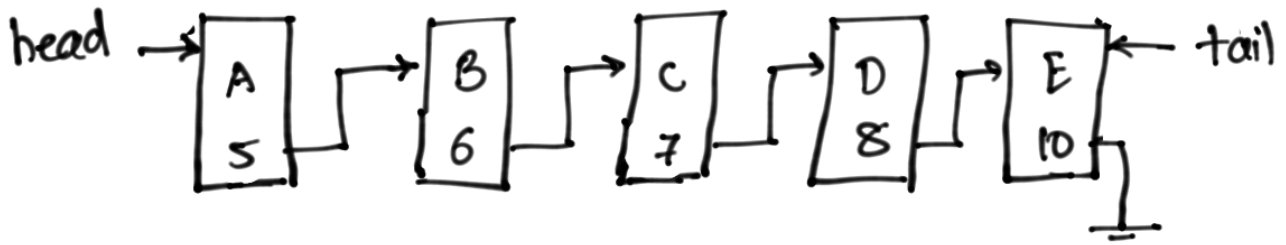
```
    }
```

```
}
```

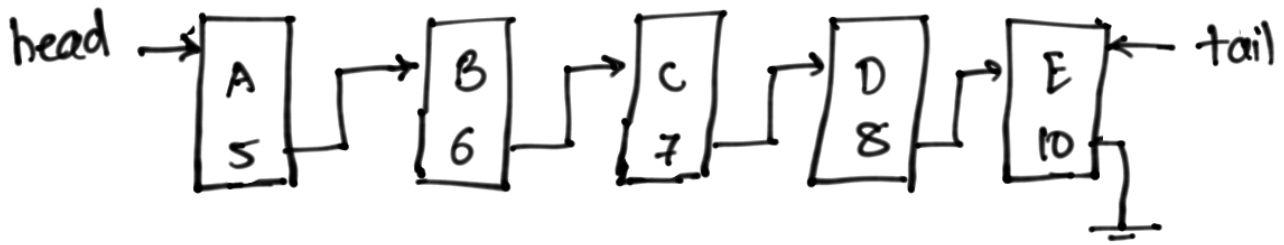
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A:  $O(1)$ .

delete Employee C

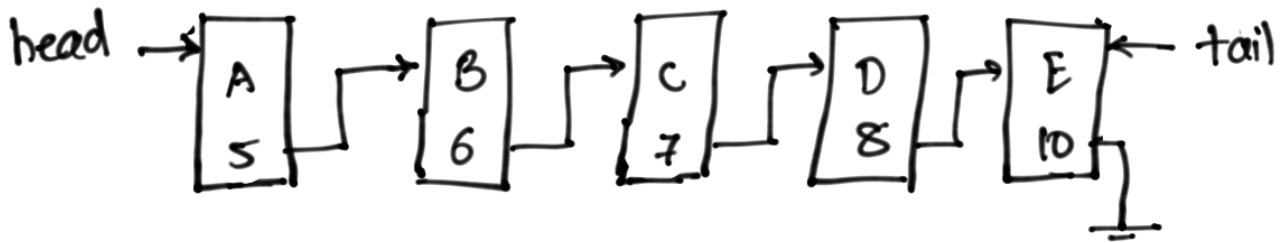


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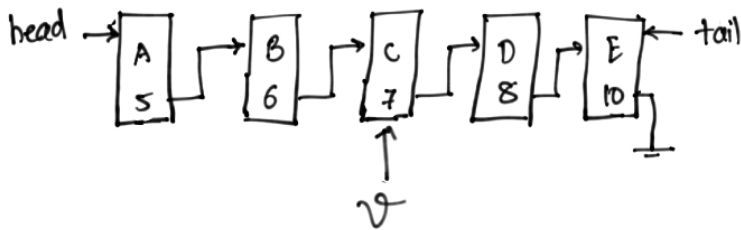
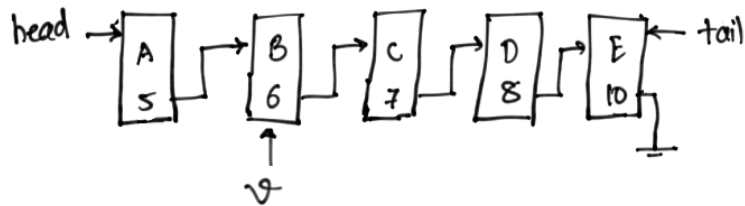
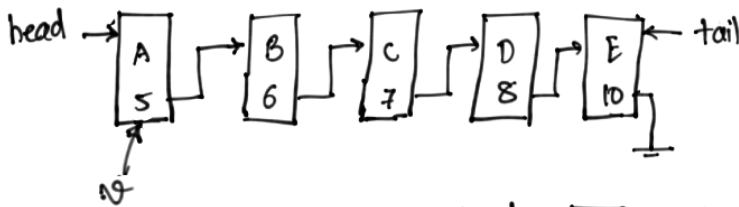
①  $v \leftarrow \text{head}$  (start with the head).

# delete Employee C



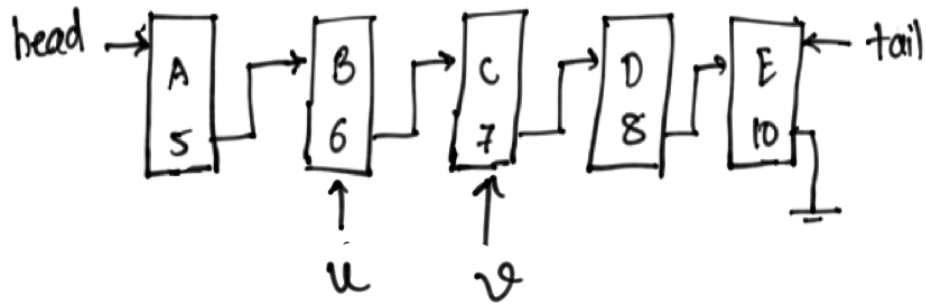
①  $v \leftarrow \text{head}$  (start with the head).

② Move through the list (using pointers) till you hit employee C

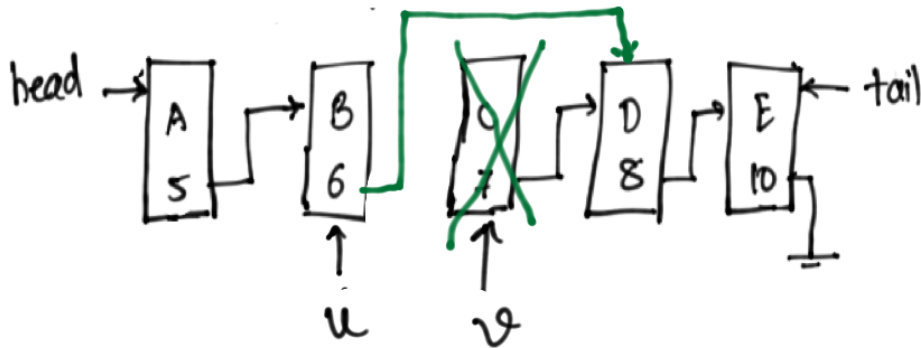




③ Maintain a ptr that follows  $v$ , say  $u$ .



④  $u.ptr \leftarrow v.ptr$   
deallocate the memory associated with  
recor



Delete(name)

{

if (head is null)

return;

Delete(name)

{

if ( head is NULL)

return;

if ( )

{

}

else

{

v ← head.ptr

u ← head;

while ( v is not NULL)

{

if (v.name = name)

{ u.ptr ← v.ptr

deallocate the memory allocated  
to record v;

}

else

{ u ← v;

v ← v.ptr

}

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if ( head is NULL)

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Q: What is the running time of  
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A  $O(n)$

# Performance of Linked List

Insert

$O(1)$

Delete

$O(n)$

Search

$O(1)$

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## Performance of Linked List

Insert	$O(1)$
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But the most important thing:  
Holy Grail for data-structure

The space taken by your data-structure should be proportional to the number of current employees in the company.

Q: When is linked list worse than arrays?

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Query: Get me  $k^{\text{th}}$  record.

Time taken by linked list = ?

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Deletion	$O(n)$
Search	$O(n)$
Get the $k^{\text{th}}$ element	$O(1)$ .

## Performance of Linked List

Insert	$O(1)$
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Get the $k^{\text{th}}$ element	$O(k)$